

Review on de Bruijn shapes in 1, 2 and 3 dimensions

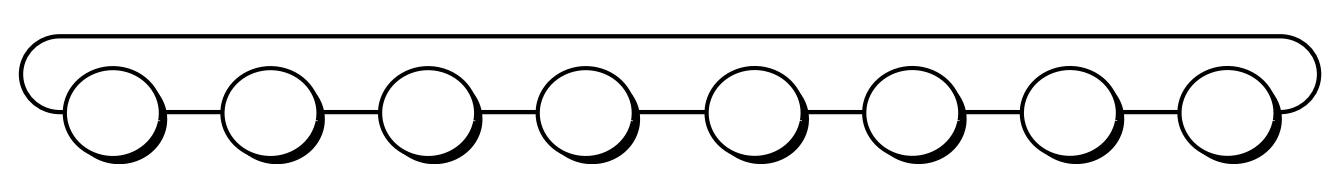
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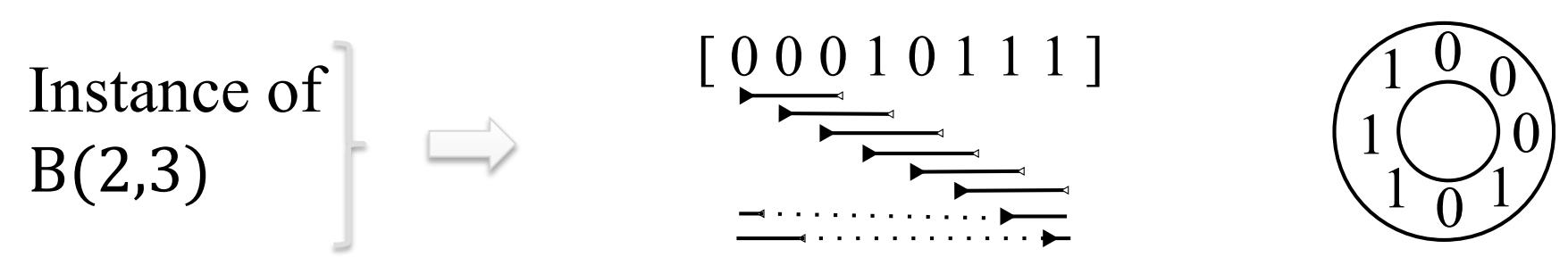
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De Bruijn sequence

- A de Bruijn sequence $B(k, n)$ is the shortest k^n -length string composed by all n -length substrings available with a k -ary alphabet containing each of those substrings just once, where $k \geq 2$ and $n \geq 1$
- It has 1 dimension and its shape is that of a toroidal array



- Reflection on a sequence gives a new instance, but rotation don't



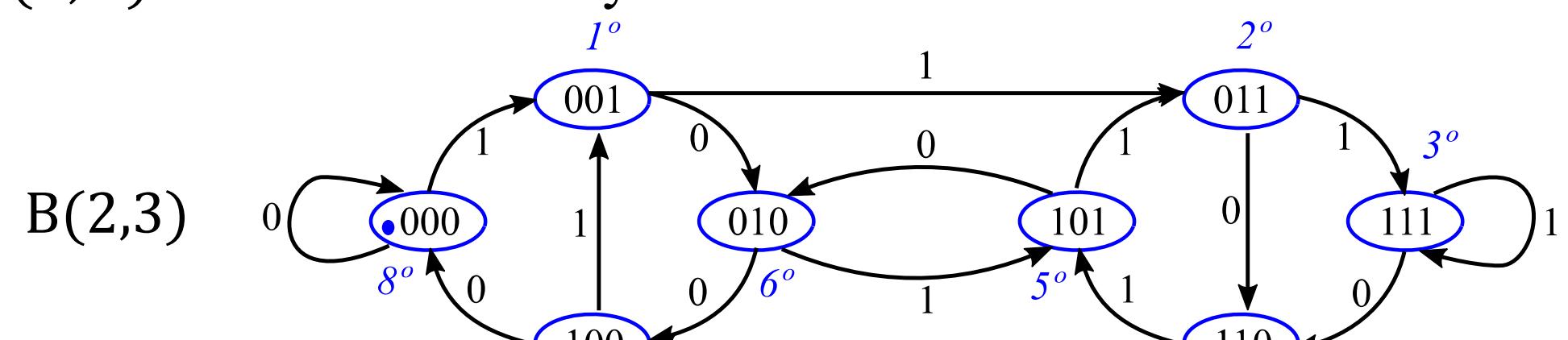
- There are $\frac{(k!)^{k^{n-1}}}{k^n}$ different sequences for given values of k and n
- They allow to precisely spot any particular substring within the string

Some instances of de Bruijn sequences

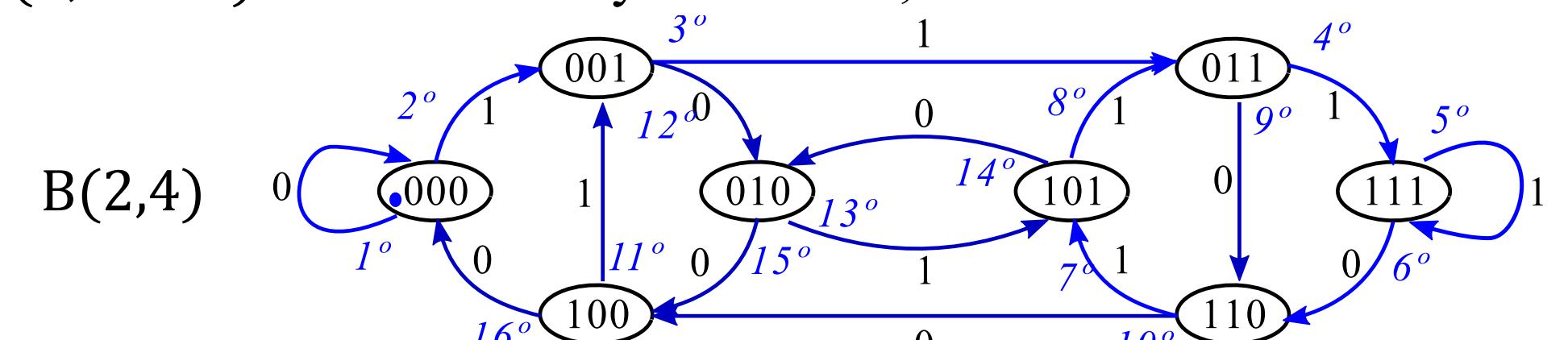
k	B(k, n)	One particular instance of that $B(k, n)$
{0, 1}	$B(2, 2)$	1100
{0, 1}	$B(2, 3)$	111010000
{0, 1}	$B(2, 4)$	11110101100010000
{0, 1}	$B(2, 5)$	11111000110111010100000100101100
{0, 1}	$B(2, 6)$	111111010101101110010110011001001100010100011000001000000
{0...2}	$B(3, 2)$	221100201
{0...2}	$B(3, 3)$	22212111000112201200201010
{0...3}	$B(4, 2)$	3322110023130120
{0...3}	$B(4, 3)$	3330300221121122020010111000312032102132013012302310331131332232
{0...4}	$B(5, 2)$	4431323303422414021120100
{0...5}	$B(6, 2)$	55445343524232251413121150403020100
{0...6}	$B(7, 2)$	665456233642446343525226155051304160314021120100
{0...7}	$B(8, 2)$	7766756557464544736353433726252423227161514131211706050403020100

De Bruijn graph

- Directed graph with k^n nodes, allowing to obtain de Bruijn sequences
- $B(k, n)$ if a Hamiltonian cycle is taken: 11101000 and its reflection 00010111



- $B(k, n+1)$ if an Eulerian cycle is done, for instance: 0111101100101000



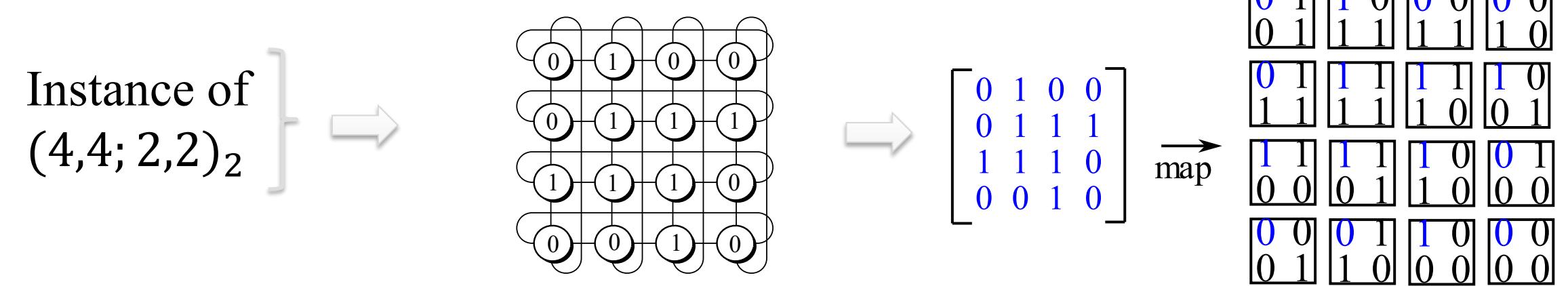
All 16 B(2,4) cycles	0000110100101111	1111001011000000	1111010010110000	0000101101001111
	0000100100101111	1111001011000000	1111010010110000	0000101001101111
	0000100100101111	1111001011000000	1111010010110000	0000101001101111
	0000100100101111	1111001011000000	1111010010110000	0000101001101111

De Bruijn Torus

- Extension of the concept of de Bruijn sequences to 2 dimensions
- A de Bruijn torus $(r_1, r_2; m_1, m_2)_k$ is the smallest $r_1 \times r_2$ matrix with all possible $m_1 \times m_2$ patterns in a k -ary alphabet just once
- Its shape is a toroidal matrix and patterns are toroidal submatrices
- A sufficient condition for a de Bruijn torus to exist:

$$r_1 \times r_2 = k^{m_1 \times m_2}$$

- They allow to exactly spot any particular 2D-pattern within the matrix
- The handle of each submatrix is its top left element



De Bruijn 3D-hypertorus

- Extension of the concept of de Bruijn torus to 3 dimensions
- A de Bruijn 3D-hypertorus $(r_1, r_2, r_3; m_1, m_2, m_3)_k$ is the smallest $r_1 \times r_2 \times r_3$ 3D-hypermatrix with all possible $m_1 \times m_2 \times m_3$ patterns in a k -ary alphabet appearing exactly once
- Its shape is a 3D-hypertoroidal matrix and patterns are 3D-hypertoroidal submatrices
- A sufficient condition for a de Bruijn 3D-hypertorus to exist:

$$r_1 \times r_2 \times r_3 = k^{m_1 \times m_2 \times m_3}$$

- They allow to spot any particular 3D-pattern within the 3D matrix
- The handle of each 3D submatrix is its front top left element
- Each handle gives the distances to the front top left corner of the 3D matrix, which is considered to be (0,0,0), such that layer 0, row 0 and column 0

Wong Algorithm

- An algorithm for binary and k -ary alphabets allowing to achieve de Bruijn sequences

o A necklace is the lexicographically earliest string within a rotation set

$$f_2(b_1 b_2 \dots b_n) = \begin{cases} b_2 b_3 \dots b_n \bar{b}_1 & \text{if } b_2 b_3 \dots b_n b_1 \text{ is a necklace;} \\ b_2 b_3 \dots b_n b_1 & \text{otherwise.} \end{cases}$$

$$f_k(a_1 a_2 \dots a_n) = \begin{cases} a_2 a_3 \dots a_n a_b & \text{if } a_1 = k-1; \\ a_2 a_3 \dots a_n (a_1 + 1) & \text{if } a_1 \neq k-1 \text{ and} \\ & a_2 a_3 \dots a_n (a_1 + 1) \text{ is a necklace;} \\ a_2 a_3 \dots a_n a_1 & \text{otherwise.} \end{cases}$$

Now: $b_1 b_2 b_3$	Test: $b_2 b_3 b_1$	Necklace?	Next: YES: $b_2 b_3 b_1$ NO: $b_2 b_3 b_1$
000	001	Yes	001
001	011	Yes	011
011	111	Yes	111
111	111	Yes	110
110	101	No	101
101	011	Yes	010
010	101	No	100
100	001	Yes	000

Some instances of de Bruijn tori

(16,32;3,3) ₂	(8,8;3,2) ₂	(4,16;3,2) ₂	(9,9;2,2) ₃	(16,16;2,2) ₄
0000001011100101100101100101111	00000101	0000011101010011	000100010	0010001030203020
00010011111010010110101000110	001001111	001100110010011011	001221221	0001020301000203
0000001011100101100101100101111	001101101011121121	111121211	0111011131213121	1011121311101213
111011000000101100101010111001	010000001	111100100110001	112002002	221201021
0000001011100101100101100101111	111101010	221201202	1011121311101213	0010001030203020
0000101111101001011010100011010	110100000	221201021	2212012232120223	02120212322322
0001001111101001011010100011010	110100000	00100010203020	002012102	0111011131213121
001100011010110111000100100100	100100001	0010001010001111	10011110	222022202
1000101001101010100000110001111	10111110	222022202	3031323331303233	0313031333233323
1111101000110100010100011000000	110100000	111100000	1011121311101213	1011121311101213
111011000000101100101010111001	111100000	111100000	221201021	02120212322322
0101000001011001010101110010000	100100000	001000000	0001020301000203	0001020301000203
0100000001011001010101110010000	101100000	101100000	1011121311101213	1011121311101213
111111000000101100101010111001	111100000	111100000	2212012232120223	2021222321202223
0101000001011001010101110010000	100100000	000100000	02120212322322	02120212322322
0100000001011001010101110010000	101100000	101100000	0111011131213121	0111011131213121

Mapping of de Bruijn 3D hypertorus (16,4,4;2,2,2)

b-cube	handle	b-cube	handle	b-cube	handle	b-cube	handle
[0 0 0]	[0 0 0]	[7,3,0]	[1 0 0]	[14,2,2]	[1 0 0]		